

## Assignment – 15

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### 1. Explain Data Encryption Standard (DES) and Rivest-Shamir-Adleman (RSA) Algorithms

**Answer:** Data Encryption Standard (DES)

The Data Encryption Standard (DES) is a symmetric-key block cipher, a foundational algorithm that played a significant role in cryptography. While no longer considered secure for modern applications due to its short key length, it's still a valuable learning tool.

**Symmetric-Key:** The same secret key is used for both encryption and decryption.

**Block Cipher:** Operates on fixed-size blocks of data (64 bits in DES).

#### **DES Process:**

1. Plaintext: The original message to be encrypted.
2. Key: A secret 56-bit key used for encryption and decryption.
3. Feistel Cipher: DES uses a Feistel structure, which involves multiple rounds of substitutions and permutations on the data being processed.
4. Ciphertext: The scrambled version of the plaintext generated by the algorithm.

#### **Limitations of DES:**

**Short Key Length (56-bit):** With increased computing power, brute-forcing the key becomes feasible, making DES vulnerable.

**Outdated:** Developed in the 1970s, DES is no longer considered secure for contemporary encryption needs.

#### **Rivest-Shamir-Adleman (RSA) Algorithm**

Rivest-Shamir-Adleman (RSA) is a widely used asymmetric-key cryptosystem, a more robust solution for modern encryption compared to DES.

**Asymmetric-Key:** Uses two mathematically linked keys - a public key for encryption and a private key for decryption.

#### **RSA Key Generation:**

1. Prime Numbers: Two large prime numbers are chosen to create a public key.
2. Public and Private Keys: Complex mathematical calculations involving these prime numbers generate a public key (shared) and a private key (kept secret).

#### **Encryption and Decryption:**

A message encrypted with the public key can only be decrypted with the corresponding private key, ensuring secure communication.

#### **Advantages of RSA:**

**Long Key Lengths:** Supports much larger key sizes (typically 2048 or 4096 bits) compared to DES, making it significantly more secure.

**Digital Signatures:** RSA can be used for digital signatures, allowing verification of the sender's identity and message integrity.

In essence:

DES is a historical algorithm, valuable for understanding encryption concepts, but not recommended for modern security needs due to its short key length.

RSA is a widely used and secure asymmetric encryption system that offers robust protection for digital communication.

## 2. Explain Diffie-Hellman Key Exchange Algorithm With an Example

**Answer:** The Diffie-Hellman (DH) key exchange is a clever cryptographic protocol that allows two parties, let's call them Alice and Bob, to establish a shared secret key over an insecure public channel without ever directly exchanging the key itself. Here's how it works:

Imagine Alice and Bob want to chat securely online, but eavesdroppers might be lurking on the internet.

### 1. Agreeing on Public Parameters:

They agree on two publicly known values beforehand:

A large prime number ( $p$ ) - This acts as the base for the mathematical calculations.

A generator ( $g$ ) - This is a number with a special property within the modulo operation with  $p$  (think of it as a mathematical starter).

### 2. Keeping Secrets Private:

Alice: Chooses a secret random number ( $a$ ) - This is her private key.

Bob: Chooses a secret random number ( $b$ ) - This is his private key.

### 3. Exchanging Public Information (Safe to Send over Insecure Channel):

Alice: Calculates  $g^a \bmod p$  (performs the exponentiation operation with modulo  $p$ ) and sends this value ( $A$ ) publicly to Bob.

Bob: Calculates  $g^b \bmod p$  and sends this value ( $B$ ) publicly to Alice.

Crucially, neither  $A$  nor  $B$  reveal their original secret numbers ( $a$  and  $b$ ).

### 4. Deriving the Shared Secret:

Alice: Receives Bob's public message ( $B$ ). She calculates  $(B)^a \bmod p$  which is essentially  $(g^b)^a \bmod p$ . Using the mathematical properties of modular exponentiation, this simplifies to  $g^{(ab)} \bmod p$ .

Bob: Receives Alice's public message ( $A$ ). He calculates  $(A)^b \bmod p$  which simplifies to  $g^{(ab)} \bmod p$  as well.

Through this magic of modular arithmetic, both Alice and Bob end up with the same shared secret key ( $g^{(ab)} \bmod p$ ) without ever revealing their private keys  $a$  and  $b$ !

Analogy with Colors (for a non-mathematical perspective):

Imagine colors represent numbers. They agree on a base color (yellow) and each keeps a secret color hidden (red and blue for Alice and Bob). They then mix their secret color with the public yellow and send the mixed color to each other. Even though they only transmit these mixed colors publicly, when they each mix their own secret color again

with the received mixed color, they both end up with the same final color, representing the shared secret.

Security of Diffie-Hellman:

The security relies on the difficulty of calculating the discrete logarithm. It's hard for an eavesdropper who only witnessed the public messages A and B (mixed colors) to figure out the original secret numbers a and b (individual secret colors) needed to derive the shared secret key.

Note: Diffie-Hellman establishes a shared secret for encryption algorithms, but doesn't provide encryption itself. It's often used in conjunction with other cryptographic techniques for secure communication.

### 3. Explain Digital Signature Algorithm (DSA) With an Example.

**Answer:** The Digital Signature Algorithm (DSA) is a cryptographic system used for generating digital signatures to verify the authenticity and integrity of digital messages. Unlike encryption which scrambles messages, DSA allows verification of the sender's identity and ensures the message hasn't been tampered with during transmission.

Here's how DSA works, along with an example:

#### 1. Key Generation:

Public and Private Keys: DSA involves a key pair - a public key (shared) for verification and a private key (kept secret) for signing messages. These keys are mathematically linked using a set of publicly known parameters (p, q, g).

Example: Let's say these are the chosen parameters:

Prime number (p) = 283

Prime number (q) = 47

Generator (g) = 60

Alice (the sender) keeps a secret random number (x) as her private key (unknown to everyone). Through a mathematical formula, she derives her public key (y) using p, q, g, and x, but x is never revealed.

#### 2. Signing a Message:

When Alice wants to sign a message (M), she first creates a hash of the message (H(M)) using a secure hashing function like SHA-256. This hash acts as a digital fingerprint of the message.

Alice then generates a random number (k) for this specific signing process (different for each message).

Using her private key (x), the message hash (H(M)), and the random number (k), Alice calculates the signature (s and r) through a specific mathematical formula.

The signature (s, r) is appended to the message (M) and sent together.

### 3. Verification Process:

Bob (the receiver) receives the message (M) and the signature (s, r).

He uses Alice's public key (y) along with the message hash (H(M)) and the received signature (s, r) to perform a verification calculation.

If the verification calculation matches, Bob can be confident that the message originated from Alice (due to her private key) and hasn't been altered (as tampering with the message would change the hash).

### Important Points:

DSA relies on the difficulty of the mathematical Discrete Logarithm Problem (DLP). It's computationally infeasible to derive Alice's private key (x) even knowing the public key (y) and the public parameters.

The random number (k) used for signing is crucial for security. Using the same k for multiple messages can compromise the system.

### Benefits of DSA:

Provides message integrity and sender authentication.

Widely used for digital signatures in applications like S/MIME emails and some digital certificates.

### Limitations:

Slower compared to some other signature algorithms like RSA.

Key sizes might need to be increased for future security needs.

## **4. Explain the Following Types of One-time Password (OTP) Algorithms with Examples:**

### **a) Time-based OTP (TOTP) b) HMAC-based OTP (HOTP)**

#### ## One-Time Password (OTP) Algorithms: TOTP vs HOTP

One-time passwords (OTPs) are an extra layer of security used during login or transactions. They provide a unique code that expires after a short period, making it more difficult for attackers to gain unauthorized access. Here's a breakdown of two common OTP algorithms:

#### a. Time-based One-Time Password (TOTP):

Concept: TOTP generates codes based on the current time. Both the user's device (phone app) and the authentication server use the same time synchronization source (usually internet time) to derive the code.

Process:

1. Shared Secret: During account setup, a secret key (in the form of a QR code or alphanumeric string) is shared between the user's device and the authentication server.

2. Time-based Window: The current time is divided into short intervals (typically 30 or 60 seconds).

3. Code Generation: A cryptographic hash function (like SHA-1) is used on the secret key combined with a counter representing the current time window. The result is truncated to generate a short numeric code (usually 6 or 8 digits).

Example:

- Imagine it's Friday, 9th June 2024, 6:43 PM. The current time window might be 6:40 PM - 6:45 PM.

- The user's device and the server have the same secret key.

- The device calculates the hash of the secret key combined with a counter representing the current time window (e.g., counter value 23 for the 23rd 30-second interval since 6:40 PM).

- Based on the hash value, a unique 6-digit code (e.g., 123456) is generated, valid only for this specific time window.

#### b. HMAC-based One-Time Password (HOTP):

Concept: HOTP generates codes based on a counter value, which is incremented with each new OTP request. Both the user's device and the server maintain synchronized counters.

Process

1. Shared Secret: Similar to TOTP, a secret key is shared between the user's device and the authentication server.

2. Counter Synchronization: Both devices maintain a counter value, which is incremented each time a new OTP is requested.

3. Code Generation: A cryptographic hash function (like HMAC-SHA1) is used on the secret key combined with the current counter value. The result is truncated to generate a short numeric code.

Example:

- The user requests an OTP for logging in.

- The user's device has a counter value of 123 (meaning 123 previous OTPs have been used).

- The device calculates the HMAC-SHA1 hash of the secret key combined with the counter value (123).

- Based on the hash value, a unique 6-digit code (e.g., 789012) is generated.

- After a successful login, both the user's device and the server increment their counter values to 124.

#### Key Differences:

Synchronization Source: TOTP relies on synchronized time, while HOTP relies on synchronized counter values.

Offline Functionality: TOTP can potentially function offline (if the device has the correct time) as long as the current time window hasn't expired. HOTP typically requires online communication with the server for counter synchronization.

Security: Both algorithms are secure when implemented correctly. However, some argue that TOTP might be slightly less vulnerable to brute-force attacks due to the moving time window.

In conclusion, both TOTP and HOTP offer robust solutions for two-factor authentication, with TOTP potentially offering some advantages in offline scenarios. The choice between them might depend on specific application requirements.